Outline

- Normal Equations Review
- Augmented Normal Equations
- Fast Methods to Solve Normal Equations
 - Levinson Durbin Algorithm
 - Schur Algorithm

Normal Equations Review

• A method to design linear forward predictor is achieved by minimizing the energy of the error function

$$a_p = \arg \min_{a_p} E\{|f(p)|^2\}$$
$$= \arg \min_{a_p} E\{f_p[n]f_p^*[n],\}$$

where

$$f_p[n] = x[n] - \hat{x}[n] = x[n] + \sum_{k=1}^p a_p[k]x[n-k]$$

or

$$f_p[n] = \sum_{k=0}^p a_p[k]x[n-k],$$

with $a_p[0] = 1$

Normal Equations Review

• The minimum of this cost function can be found by taking the derivative with respect to $a'_p s$ yielding

$$\sum_{k=0}^{p} a_p[k] R_x(l-k) = 0, \quad l = 1, 2, \dots, p$$

- This is called the normal equation
- When the solution of this equation is substituted into the definition of error function, we obtain the minimum error

$$E_p^f = \sum_{k=1}^p a_p[k] R_x(-k)$$

Augmented Normal Equations

• If we combine the normal equations with the minimum error resulting from the optimum $a'_p s$, we obtain the augmented normal equations

$$\sum_{k=0}^{p} a_p[k] R_x(l-k) = \left\{ \begin{array}{cc} E_p^f, & l=0\\ 0, & l=1,2,\dots,p \end{array} \right\}$$

Fast Methods to Solve Normal Equations

- Normal equations are linear equations in $a'_p s$ and can be solved easily given the second order statistics of the signal $R_x(\tau)$
- This requirement actually explains how a future value can be predicted. We simply know the statistics (related to the likelihood of a certain value)
- Although the solution is simply the solution to a linear set of equations, computation can be high when the signal length p is high
- Therefore, fast methods have been proposed to solve linear equations utilizing the special structure of the autocorrelation function $R_x(\tau)$
- We will study two of these methods
 - Levinson-Durbin algorithm
 - Schur Algorithm

Levinson-Durbin Algorithm

- This is a recursive algorithm that utilizes the toeplitz and symmetry properties of the autocorrelation function
- Let us write the normal equations in the matrix form:

$$\Gamma_p \boldsymbol{a}_p = -\boldsymbol{r}_x$$

where a_p is the vector of uknown coefficients, and Γ_p is the matrix of the autocorrelation R_x , r_x is a vector obtained from elements of $R_x(.)$

$$\Gamma_{p} = \begin{bmatrix} R_{x}(0) & R_{x}(-1) & \dots & R_{x}(1-p) \\ R_{x}(1) & R_{x}(0) & \dots & R_{x}(2-p) \\ \dots & \dots & \dots & \dots \\ R_{x}(p-1) & R_{x}(p-2) & \dots & R_{x}(0) \end{bmatrix}$$

• Let us rewrite Γ_p using the symmetry property of the autocorrelation function $R_x(\tau) = R_x^*(-\tau)$

$$\Gamma_{p} = \begin{bmatrix} R_{x}(0) & R_{x}^{*}(1) & \dots & R_{x}^{*}(p-1) \\ R_{x}(1) & R_{x}(0) & \dots & R_{x}^{*}(p-2) \\ \dots & \dots & \dots & \dots \\ R_{x}(p-1) & R_{x}(p-2) & \dots & R_{x}(0) \end{bmatrix}$$

- This matrix is Hermitian: $\Gamma_p^{\mathrm{T}} = \Gamma_p^*$ and it is to eplitz: $\Gamma_p(i,j) = \Gamma_p(i-j)$
- Normally a matrix inversion takes time $O(p^3)$, but using these properties we will be able to solve the system more efficiently

• Let us start by calculating $a_1(1)$

$$a_1(1) = -\frac{R_x(1)}{R_x(0)}$$

with the minimum mean squared error (MMSE)

$$E_1^f = R_x(0) + a_1(1)R_x(-1)$$

= $R_x(0)(1 - |a_1(1)|^2)$

• Now let us calculate the second order predictor coefficients

$$a_2(1)R_x(0) + a_2(2)R_x^*(1) = -R_x(1)$$

$$a_2(1)R_x(1) + a_2(2)R_x(0) = -R_x(2)$$

resulting in the solution

$$a_2(2) = -\frac{R_x(2) + a_1(1)R_x(1)}{E_1^f}$$

$$a_2(1) = a_1(1) + a_2(2)a_1^*(1)$$

- Note that the values a_2 's are calculated in terms of the previous step
- Continuing in this manner we can obtain all coefficients

$$\boldsymbol{a}_{m} = \begin{bmatrix} a_{m}(1) \\ a_{m}(2) \\ \vdots \\ \vdots \\ \vdots \\ a_{m}(m) \end{bmatrix} = \begin{bmatrix} \boldsymbol{a}_{m-1} \\ \vdots \\ \vdots \\ \vdots \\ 0 \end{bmatrix} + \begin{bmatrix} \boldsymbol{d}_{m-1} \\ \vdots \\ \vdots \\ \vdots \\ K_{m} \end{bmatrix}$$

• What are the values of d_m 's and K_m ?

• Let us partition Γ_m as follows

$$\Gamma_m = \left[egin{array}{ccc} \Gamma_{m-1} & m{r}_{m-1}^{r*} \ m{r}_{m-1}^{\mathrm{rT}} & R_x(0) \end{array}
ight],$$

where $\mathbf{r}_{m-1}^r = [R_x(m-1)R_x(m-2), \dots, R_x(1)]$

• The normal equations can now be written as

$$\left[\begin{array}{cc} \Gamma_{m-1} & \boldsymbol{r}_{m-1}^{r*} \\ \boldsymbol{r}_{m-1}^{r\mathrm{T}} & R_{x}(0) \end{array}\right] \left\{ \left[\begin{array}{c} \boldsymbol{a}_{m-1} \\ 0 \end{array}\right] + \left[\begin{array}{c} \boldsymbol{d}_{m-1} \\ K_{m} \end{array}\right] \right\} = - \left[\begin{array}{c} \boldsymbol{r}_{m-1} \\ R_{x}(m) \end{array}\right]$$

Levinson-Durbin Algorithm

• We have two equations from two rows:

$$\Gamma_{m-1} a_{m-1} + \Gamma_{m-1} d_{m-1} + K_m r_{m-1}^{r*} = -r_{m-1}$$

$$r_{m-1}^{rT} a_{m-1} + r_{m-1}^{rT} d_{m-1} + K_m R_x(0) = -R_x[m]$$

• Solving the first row of this block form yields

$$d_{m-1} = -K_m \Gamma_{m-1}^{-1} r_{m-1}^{r*}$$

= $K_m a_{m-1}^{r*}$ (*)

where r again denotes reversing of the element order since $\Gamma_{m-1}a_{m-1} = -r_{m-1}.$

• Now using second row equation and Eq. (*) we obtain

$$K_m[R_x(0) + \boldsymbol{r}_{m-1}^{\mathrm{rT}} \boldsymbol{a}_{m-1}^*] + \boldsymbol{r}_{m-1}^{\mathrm{T}} \boldsymbol{a}_{m-1} = -R_x(m)$$

• Hence, we obtain K_m

$$K_m = -\frac{R_x(m) + \boldsymbol{r}_{m-1}^{\mathrm{rT}} \boldsymbol{a}_{m-1}}{R_x(0) + \boldsymbol{r}_{m-1}^{\mathrm{rT}} \boldsymbol{a}_{m-1}^*}$$

• We have obtained K_m and d_m , the whole recursive set of equations are

$$a_{m}(m) = K_{m} = -\frac{R_{x}(m) + \boldsymbol{r}_{m-1}^{\mathrm{T}} \boldsymbol{a}_{m-1}}{R_{x}(0) + \boldsymbol{r}_{m-1}^{\mathrm{T}} \boldsymbol{a}_{m-1}^{*}}$$

$$a_{m}(k) = a_{m-1}(k) + K_{m} a_{m-1}^{r*}(m-k)$$

$$k = 1, \dots, m-1$$

• The K_m 's are the reflection coefficients in the lattice structure

• The expression for the MMSE is

$$E_m^f = R_x(0) + \sum_{k=1}^m a_m(k) R_x(-k)$$

= $R_x(0) + \sum_{k=1}^m [a_{m-1}(k) + a_m(m) a_{m-1}^*(m-k)] R_x(-k)$
= $E_{m-1}^f (1 - |K_m|^2)$

- The error is monotonically decreasing, makes sense since as new data comes in we should have improving performance
- Computational cost: for each stage we have O(m) multiplications resulting in total $1 + 2 + \ldots + p = p(p+1)/2$ that is $O(p^2)$ operations
- Regular inversion $O(p^3)$ for arbitrary matrix

Schur Algorithm

• Consider the following

$$\alpha_0(z) = \frac{R_x(1)z^{-1} + R_x(2)z^{-2} + \ldots + R_x(p)z^{-p}}{R_x(0) + R_x(1)z^{-1} + R_x(2)z^2 + \ldots + R_x(p)z^{-p}}$$

and the following that will be calculated recursively

$$\alpha_m(z) = \frac{\alpha_{m-1}(z) - \alpha_{m-1}(\infty)}{z^{-1}[1 - \alpha_{m-1}^*(\infty)\alpha_{m-1}(z)]}$$

• Since $\alpha_0(\infty) = 0$ we can obtain

$$\alpha_1(z) = \alpha_0(z)/z^{-1} = \frac{R_x(1) + R_x(2)z^{-1} + \ldots + R_x(p)z^{-p+1}}{R_x(0) + R_x(1)z^{-1} + R_x(2)z^2 + \ldots + R_x(p)z^{-p}}$$

- Now with these definitions, we have $\alpha_1(\infty) = R_x(1)/R_x(0)$ which is equal to $-K_1$
- The next step in the recursion

$$\alpha_2(\infty) = \frac{R_x(2) + K_1 R_x(1)}{R_x(0)(1 - |K_1|^2)}$$

which is equal to K_2

- That is we have the relation $\alpha_m(\infty) = -K_m$
- Calculating $\alpha_m(\infty)$'s is equivalent to solving normal equations

• Let us write

$$\alpha_m(z) = \frac{P_m(z)}{Q_m(z)}$$

where

$$P_0(z) = R_x(1)z^{-1} + R_x(2)z^{-2} + \ldots + R_x(p)z^{-p}$$
$$Q_0(z) = R_x(0) + R_x(1)z^{-1} + R_x(2)z^{-2} + \ldots + R_x(p)z^{-p}$$

• Consider the following block equations for recursive creation of $P_m(z)$ and $Q_m(z)$

$$\begin{bmatrix} P_m(z) \\ Q_m(z) \end{bmatrix} = \begin{bmatrix} 1 & K_{m-1} \\ K_{m-1}^* z^{-1} & z^{-1} \end{bmatrix} \begin{bmatrix} P_{m-1}(z) \\ Q_{m-1}(z) \end{bmatrix}$$

• Let us check the recursive relation

 $P_1(z) = P_0(z)$ $Q_1(z) = z^{-1}Q_0(z)$

• Next step

$$P_{2}(z) = P_{1}(z) + K_{1}Q_{1}(z)$$

= $[R_{x}(2) + K_{1}Q_{1}(z)]$
 $Q_{2}(z) = z^{-1}[Q_{1}(z) + K_{1}^{*}P_{1}(z)]$

- We can see that $P_2(z)/Q_2(z)$ is equivalent to the definition of $\alpha_2(z)$
- Recursive applications show that this is true for all m

- The Steps of Schur Algorithm can systematically performed as the following
 - Create a matrix

$$G = \begin{bmatrix} 0 & R_x(1) & r_x(2) & \dots & R_x(p) \\ R_x(0) & R_x(1) & r_x(2) & \dots & R_x(p) \end{bmatrix}$$

 Shift second row to right once with padding zeros for the new elements, the negative ratio of the second column is the reflection coefficient

- And then
 - Multiply this shifted matrix by

$$\begin{bmatrix} 1 & K_m \\ K_m^* & 1 \end{bmatrix}$$

for (m = 1, 2, ...)

- Shift to right by once again, the negative ratio of the second column is ${\cal K}_2$
- Continue untill all K's are calculated